Multicore Design Considerations

Multicore: The Forefront of Computing Technology

"We're not going to have faster processors. Instead, making software run faster in the future will mean using parallel programming techniques. This will be a huge shift." -- Katherine Yelick, Lawrence Berkeley National Laboratory from The Economist: Parallel Bars

- Multicore is a term associated with parallel processing, which refers to the use of simultaneous processors to execute an application or multiple computational threads.
- Parallel programming/processing can be implemented on TI's KeyStone multicore architecture.

Parallel Processing

- Parallel processing divides big applications into smaller applications and distributes tasks across multiple cores.
- The goal is to speed up processing of a computationallyintensive applications.
- Characteristics of computationally-intensive applications:
 - Large amount of data to process
 - Complex algorithms require many computations
- Goals of task partitioning
 - Computational load balancing evenly divides effort among all available cores
 - Minimizes contention of system resources
 - Memory (DDR, shared L2)
 - Transport (Teranet, peripherals)

Parallel Processing: Use Cases

- Network gateway, speech/voice processing
 - Typically hundreds or thousands of channels
 - Each channel consumes about 30 MIPS
- Large, complex, floating point FFT (1M)
- Multiple-size, short FFTs
- Video processing
 - Slice-based encoder
 - Video transcoder (low quality)
 - High-quality decoder

Parallel Processing: Use Cases

- Medical imaging
 - Filtering > reconstruction > post filtering
 - Edge detection
- LTE channel excluding turbo decoder/encoder
 - Two cores uplink
 - Two cores downlink
- LTE channel including turbo decoder
 - Equal to the performance of 30 cores
 - Each core works on a package of bits
- Scientific processing
 - Large complex matrix manipulations
 - Use Case: Oil exploration

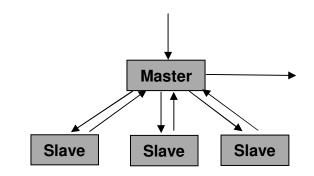
Parallel Processing: Control Models

Master Slave Model

- Multiple speech processing
- Variable-size, short FFT
- Video encoder slice processing
- VLFFT

Data Flow Model

- High quality video encoder
- Video decoder
- Video transcoder
- LTE physical layer





Parallel Processing: Partitioning Considerations

Function driven

- Large tasks are divided into function blocks
- Function blocks are assigned to each core
- The output of one core is the input of the next core
- Use cases: H.264 high quality encoding and decoding, LTE

Data driven

- Large data sets are divided into smaller data sets
- All cores perform the same process on different blocks of data
- Use cases: image processing, multi-channel speech processing, sliced-based encoder

Parallel Processing: System Recommendations

- Ability to perform many operations
 - Fixed-point AND floating-point processing
 - SIMD instruction, multicore architecture
- Ability to communicate with the external world
 - Fast two-way peripherals that support high bit-rate traffic
 - Fast response to external events
- Ability to address large external memory
 - Fast and efficient save and retrieve methods
 - Transparent resource sharing between cores
- Efficient communication between cores
 - Synchronization
 - Messaging
 - Data sharing

Parallel Processing: Recommended Tools

- Easy-to-use IDE (Integrated Development Environment)
 - Advanced debug features (system trace, CP tracer)
 - Simultaneous, core-specific debug monitoring
- Real-time operating system (e.g., SYS/BIOS)
- Multicore software development kit
 - Standard APIs simplifies programming
 - Layered abstraction hides physical details from the application
- System optimized capabilities
 - Full-featured compiler, optimizer, linker
 - Third-party support

Example: High Def 1080i60 Video H264 Encoder

- A short introduction to video encoding
 - Pixel format
 - Macroblocks
- Performance numbers and limitations
 - Motion estimation
 - Encoding
 - Entropy encoder
 - Reconstruction
 - Data in and out of the system
 - DDR bandwidth
 - Synchronization, data movement
 - System architecture

Macroblock and Pixel Data

RGB and YUV

$$\begin{bmatrix} Y' \\ U \\ V \end{bmatrix} = \begin{bmatrix} 0.299 & 0.587 & 0.114 \\ -0.14713 & -0.28886 & 0.436 \\ 0.615 & -0.51499 & -0.10001 \end{bmatrix} \begin{bmatrix} R \\ G \\ B \end{bmatrix}$$
$$\begin{bmatrix} R \\ G \\ B \end{bmatrix} = \begin{bmatrix} 1 & 0 & 1.13983 \\ 1 & -0.39465 & -0.58060 \\ 1 & 2.03211 & 0 \end{bmatrix} \begin{bmatrix} Y' \\ U \\ V \end{bmatrix}$$

4:4:4 and 4:2:0 format

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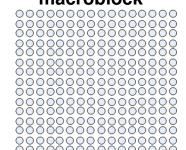
4:4:4

4:2:0

- -- Pixel with only Y value
- -- Pixel with only Cr and Cb values
- -- Pixel with Y, Cr, and Cb values

Typically 8-bit values (10, 12, 14)

Macroblock = 16x16 pixels



Video Encoder Flow (per Macroblock)

Coder	Width	Height	Frames/Second	MCycles/Second
D1(NTSC)	720	480	30	660
D1 (PAL)	720	576	25	660
720P30	1280	720	30	1850
1080i	1920	1080 (1088)	60 fields	3450

Motion estimation	
Intra prediction and Motion Compensation	
\	
Integer transform	
•	
Quantization	Out Entropy Encoder
₩,	(Cabac or
Inverse Quantization	CAVLC)
 	
Inverse Integer transform And reconstruction	
reconstruction	
De-blocking Filter and re-	

Module	Percentage	Approximate MIPS	Number of Cores
		(1080i)/Second	
Motion Estimation	~50%	1750	2
IP, MC, Transform,	~12.5%	437.7	0.5
Quantization			
Entropy Encoder	~25%	875	1
IT, IQ and	~12.5%	437.5	0.5
Reconstruction			

Video Coding Algorithm Limitations

Motion estimation

- Depends on the reconstruction of previous (and future) frames
- Shortcuts can be performed (e.g., first row of frame N does not need last row of frame N-1).
- Intra-prediction
 - Depends on the macroblock above and to the left
 - Must be done consecutively or encoding efficiency is lost (i.e., lower quality for the same number of bits)
- Entropy encoding (CABAC, CAVLC)
 - Must be processed in the macroblock order
 - Each frame is independent of other frames.

How Many Channels Can One C6678 Process?

- Looks like two channels;
 Each one uses four cores.
 - Two cores for motion estimation
 - One core for entropy encoding
 - One core for everything else
- What other resources are needed?
 - Streaming data in and out of the system
 - Store and load data to and from DDR
 - Internal bus bandwidth
 - DMA availability
 - Synchronization between cores, especially if trying to minimize delay

What are the System Input Requirements?

- Stream data in and out of the system:
 - Raw data: 1920 * 1080 * 1.5 = 3,110,400 bytes per frame
 = 24.883200 bits per frame (~25M bits per frame)
 - At 30 frames per second, the input is 750 Mbps
 - NOTE: The order of raw data for a frame is Y component first, followed by U and V
- 750 Mbps input requires one of the following:
 - One SRIO lane (5 Gbps raw, about 3.5 Gbps of payload),
 - One PCle lane (5 Gbps raw)
 - NOTE: KeyStone devices provide four SRIO lanes and two PCIe lanes
- Compressed data (e.g., 10 to 20 Mbps) can use SGMII (10M/100M/1G) or SRIO or PCIe.

How Many Accesses to the DDR?

- For purposes of this example, only consider frame-size accesses.
- All other accesses (ME vectors, parameters, compressed data, etc.) are negligible.
- Requirements for processing a single frame:
 - Retrieving data from peripheral to DDR 25M bits = 3.125MB
 - Motion estimation phase reads the current frame (only Y) and older Y component of reconstruction frame(s).
 - A good ME algorithm may read up to 6x older frame(s).
 - 7 * 1920 * 1088 = ~ 15M Bytes
 - Encoding phase reads the current frame and one old frame. The total size is about 6.25 MB.
 - Reconstruction phase reads one frame and writes one frame. So the total bandwidth is 6.25 MB.
 - Frame compression before or after the entropy encoder is negligible.
 - Total DDR access for a single frame is less than 32 MB.

How Does This Access Avoid Contention?

- Total DDR access for a single frame is less than 32 MB.
- The total DDR access for 30 frames per second (60 fields) is less than 32 * 30 = 960 MBps.
- The DDR3 raw bandwidth is more than 10 GBps (1333 MHz clock and 64 bits). 10% utilization reduces contention possibilities.
- DDR3 DMA uses TeraNet with clock/3 and 128 bits.
 TeraNet bandwidth is 400 MHz * 16B = 6.4 GBps.

KeyStone SoC Architecture Resources

- 10 EDMA transfer controllers with 144 EDMA channels and 1152 PaRAM (parameter blocks)
 - The EDMA scheme must be designed by the user.
 - The LLD provides easy EDMA usage.
- In addition, Navigator has its own PKTDMA for each master.
- Data in and out of the system (SRIO, PCIe or SGMII) is done using the Navigator.
- All synchronization between cores and moving pointers to data between cores is done using the Navigator.
- IPC provides easy access to the Navigator.

Conclusion

 Two H264 high-quality 1080i encoders can be processed on a single TMS320C6678

System Architecture

